Applied Operating Systems User Perspective

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Introduction

- The OS offers its services to user programs through the system call interface.
- Often there is an additional layer between user programs and the kernel.
- This function is usually performed by the C library in Unix systems.
- Before we deal with system functions we will look at the Unix shell.

Unix Shells

- The shell is a user program.
- It works as a command interpreter.
- When a user types the name of an executable, the shell creates a process (a child) to execute the program.
- There are many types of shells, sh, csh, bash
- Most Unix executables read from standard input and write to standard output

- When a user logs in, the shell starts by typing the prompt which tells the user it is waiting for commands.
- The prompt is usually some symbol like the dollar sign or a string followed by such symbol.
- example

```
$
$ date
Thu Sep 23 18:08:44 EEST 2004
$
```

Unix Utilities

- Unix system usually came with hundreds of utility programs.
- Each one does one thing only.
- All of them use the standard input/output.
- By combining them, complicated commands can be executed.
- The shell uses system functions to redirect the output of one executable to be the input of another
- A key concept is output redirection and pipes

Redirection

- The shell interprets the symbols > and < as input/output redirection.</p>
- The > symbol redirects output. Example
 - \$ date > file
- Redirects the output of the date command into the file file.
- Similarly the < symbol redirects input. sort is a program to sort the input in alphabetical order. Example
 - \$ sort < file1 > file2
- Will read the content of file1, sort it and store the output in file2
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Pipes

- The symbol for a pipe is |
- The output of one program can be connected to the input of another using a pipe.

cat file1 file2 file3 | sort >lpr

- The cat program reads the file and prints it to standard output.
- **•** The **Ipr** file is the printer device.
- In Unix almost all devices have a file interface.
- In the above example the output of cat is connected to the input of sort and the output of sort is redirected to the printer device.

How Does The Shell Work?

The main job of the shell is

- Execute programs on behalf of the user.
- Optionally pass appropriate parameters to the program.
- Redirect input/output if needed.
- Create pipes to connect the input/output of programs.
- All the above are done using function calls provided by the system.
- The function are typically wrapper function for system calls provided by the OS.

Creating Processes

- Unix processes are created using the fork() function call.
- fork creates a child process of the current process.
- The child process is a copy of the parent process.
- The fork() function call returns 0 to the child and the process id (PID) of the child to the parent.
- The parent of all processes is the init process.

```
#include <sys/types.h>
#include <unistd.h>
```

```
int main()
{ pid_t pid;
```

```
pid=fork();
```

```
if(pid==0)
```

printf("Child process\n");

else

printf("parent process, child id=%d\n",pid);

Child Memory

- The child's memory image is a copy of the parent's.
- All the child variables are inherited from the parent and have the same value up to the fork() call.
- Since the child is a copy of the parent any change made after the **fork()** call in one of them is independent of the other.

Consider the following code and its output

```
int main() {
```

```
pid_t pid; int var=1;
```

```
var++;pid=fork();
```

```
if(pid==0)
```

printf("child &var=%x var=%d",&var,var); else

printf("parent &var=%x var=%d", &var, var); }

output

parent &var=bffffd40 var=2
child &var=bffffd40 var=2

```
int main() {
pid_t pid; int var=1;
var++;pid=fork();
if(pid==0){
 var++;
 printf("child &var=%x var=%d", &var, var);
 }
else
  printf("parent &var=%x var=%d", &var, var); }
```

output

parent &var=bffffd40 var=2
child &var=bffffd40 var=3

Who Finishes First?

- Both parent and child proceed with execution from the point of the **fork**.
- One cannot tell which one finishes first.
- It depends on the amount of work each has to do.
- If parent needs to wait for the child to terminate we should use the wait system call.

```
int main() {
pid_t pid;
int status;
pid=fork();
if(pid==0)
  printf("child\n");
else{
   wait(&status);/* parent hangs
    until child is done */
   printf("child is done\n");
   }
```

The Exec Calls

- **fork** creates a copy of the calling process.
- Many applications require the child to execute different code from the parent.
- The exec family of functions provide a way for a process to execute arbitrary code.
- The new image completely replaces the old image.
- This is the reason why no code after the exec call is executed.

The Execl family

- The path is the name of the executable with the full path.
- file is the name of the executable.
- envp[] is an array of strings holding variable-value pairs.

- 1 int main(){
- 2 if(execl("/usr/bin/ls","ls","-l",0)<0) {</pre>
- 3 printf("execl error");

```
4 exit(1);
```

}

- If execl is successful, line 3 is never executed.
- The whole executable is replaced by /usr/bin/ls.

The Argv Array

- The argv parameter passed as argument to the main function contains the command line arguments.
- argv[0] is always the executable name, followed by the other parameters in order of appearance.
- All the exec functions allow for the passing of the argv parameter.
- In the previous example: argv[0]="ls", argv[1]="-l".
- Note that the list must terminate with a NULL.

Environment Variables

- Unix uses many variable-value pairs called environment variables.
- Many utilities use the value of theses variables.
- One particularly important variable is the PATH variable.
- The PATH contains a list of directories to be searched for executables.
- By using the PATH variable one doesn't need to specify the absolute path of the executables.

The Execv Family

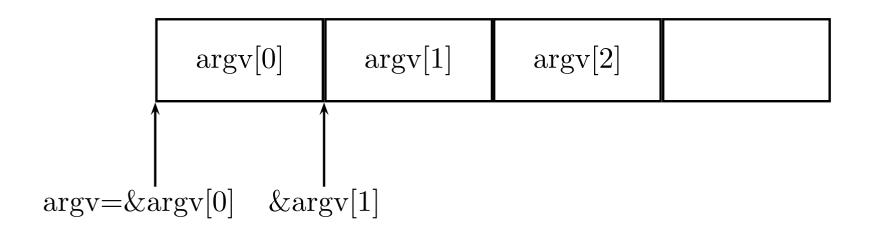
- The execv family takes the arguments for the executable as an array instead of a list.
 - int execv(char *path, char *argv[]);
 - int execvp(char *file, char *argv[]);

- If the parameter is **path** the full path needs to be specified.
- If the parameter is file the PATH variable is used to search for the executable.
- If the execve function is used one can specify the environment for the executable.

```
int main(int argc, char *argv[]) {
pid_t pid;
pid=fork();
if(pid==0){
  execvp(argv[1], & argv[1]);
  printf("error execvp");
else
wait(&status);
}
```

The above example executes any program passed on the command line along with its arguments.

Why Does It Work?

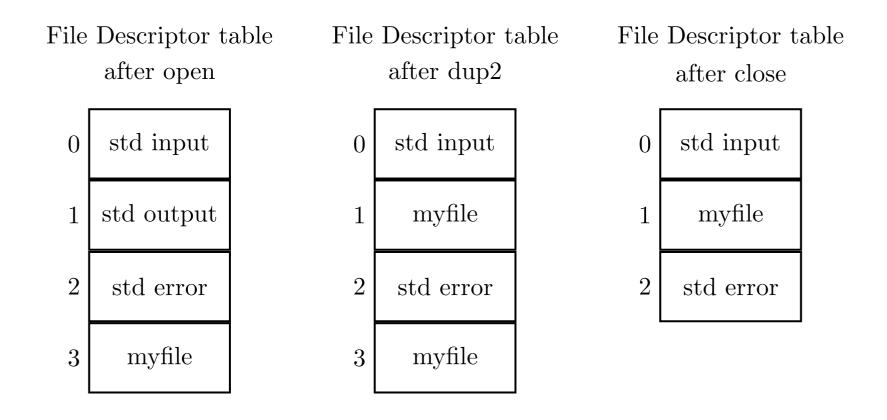


Redirection

- We have already seen that the shell can redirect the input/output of a program to a file.
- The shell does this by using the dup2 system call.
- The dup2 system call redirects the input/output of one file descriptor to another.
- Therefore to redirect output to file myfile
 - 1. Open *myfile*.
 - 2. use **dup2** to replace standard output by the descriptor of *myfile*.

```
int main() {
  int fd;
  mode_t mode=S_IRUSR|S_IWUSR|S_IRGRP|S_IROTH;
  fd=open("myfile",O_WRONLY|O_CREAT,mode);
  dup2(fd,1);
  close(fd);
  printf("test");
  }
```

- In the above example the string "test" is written to *myfile*.
- Anything written to standard output is automatically redirected to the file *myfile*.



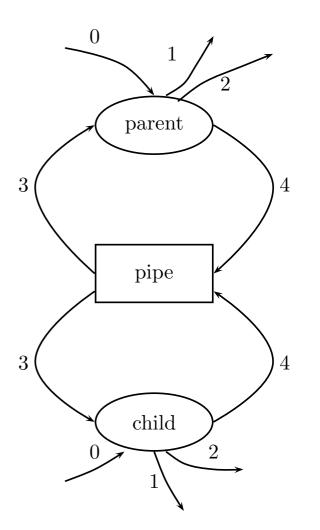
Pipes

- A pipe is a communication buffer that connects the standard output of one program to the standard input of another.
- A pipe has no external or permanent name.
- Thus it is used only by the process that created it and by its descendents.
- The prototype for the system call is int pipe(int fildes[2]);
- Data written to *fildes[1]* is read from *fildes[0]* in a FIFO fashion.

Example: ls -flsort

```
int main() {
int fd[2];pid_t pid;
pipe(fd);
pid=fork();
if(pid==0){
   dup2(fd[1],1);close(fd[0]);close(fd[1]);
  execl("/usr/bin/ls","ls","-l",NULL);
   }
else{
  dup2(fd[0],0);close(fd[0]);close(fd[1]);
  execl("/usr/bin/sort", "sort", NULL);
```

File Descriptors After pipe



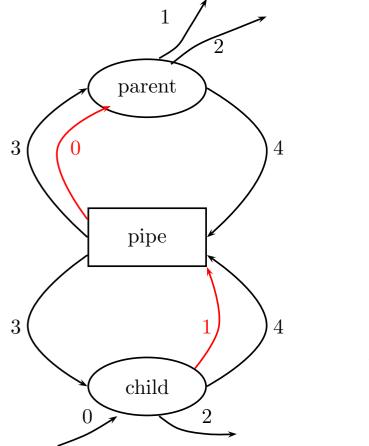
Parent file descriptor table		
0	std input	
1	std output	
2	std error	
3	pipe read	
4	pipe write	

d	Child file escriptor tabl	le
0	std input	
1	std output	
2	std error	

- 3 pipe read
- 4 pipe write

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File Descriptors After dup



Parent file descriptor table		
0	pipe read	
1	std output	
2	std error	
3	pipe read	
4	pipe write	

Child file descriptor table

0	std input
1	pipe write
2	std error
3	pipe read
4	pipe write

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Mini Shell

```
int main(int argc, char **argv) {
pid_t pid; int status,nc;
char *buf; char **args;
 buf=(char *)malloc(1024);
  while (1) {
 printf("myShell$");fflush(stdout);
  nc=read(0,buf,1024); args=parse(buf);
 buf[nc-1]=0;pid=fork();
  if(pid==0){
    execvp(args[0], args);
    printf("execvp failed\n");
  }
  else {
   wait(&status);free(args);
```

Parsing Command Line

```
char ** parse(char *buf)
{
  int count=0;char **argv;
  argv=(char **)malloc(1024);
  argv[count]=buf;
  while(*buf!=0) {
      if(*buf==' ') {
     *buf=0;count++;
     argv[count]=buf+1;
    }
      buf++;
    }
    argv[count+1]=0;
  return argv;
```